

# BERG

## TECHNICAL WHITE PAPER

The global leading privacy protecting platform  
Making decentralized applications truly Secure, Private and Stable



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**BERG**

## **CHAPTER I**

### **ABSTRACT**

The Internet has greatly enhanced the efficiency of information dissemination, which benefits Human society; on the other hand, lack of privacy becomes more of a serious problem. Blockchain is considered a great tool to protect privacy. However, since all the transactions are recorded on the public blockchain, once the identity of the wallet holder gets uncovered, this loss of privacy is irreversible. The scenario leads to a more serious problem than the privacy disclosure of the Internet. For this reason, cryptographers and top technical experts in the blockchain industry have made relentless efforts to resolve the issue. Several teams in the industry have developed special cryptocurrencies to protect privacy, which are called "anonymous currencies". Some of the best-known anonymous currencies are Zcash (ZEC), Monero (XMR), and Dash. These cryptocurrencies with a certain degree of privacy protection, have obtained high market values based on the vast demand and have been ranked among the world's top 20 cryptocurrencies for a long time; thus, indicating a strong demand for privacy protection in the blockchain industry.

Smart contract is a computer protocol designed to distribute, verify or execute contracts in an information-based way. Turing complete smart contract system on the blockchain allows developers to write any complicated contract that lives on the blockchain and can be executed on the blockchain. Developers can use smart contract development language to produce functions such as custom token, financial derivatives, identity system, and decentralized organization, therefore, greatly expanding the application scope of the blockchain system. Smart contract is one of the foundational bases of the Internet of Value. The current shortcoming is that none of the blockchain systems support encryption and privacy protection of smart contracts. The existing use scenarios of privacy protection mechanisms are greatly reduced due to the technical limitation. Blockchain 1.0 technology originated from Bitcoin invented by Satoshi Nakamoto, has created a new paradigm. With the advent of Ethereum – blockchain 2.0, the invention of smart contracts makes the blockchain technology accessible, and the Distributed Applications (DAPPs) based on the blockchain technology more feasible, allowing blockchain technology to be applicable to more industries. Zcash and Monero which do not support smart contracts are privacy protection scheme 1.0; privacy protection scheme 2.0 that supports smart contracts is expected to be implemented in more industries and application scenarios.

There is no doubt about the high technical threshold required for developing anonymous cryptocurrencies that support smart contracts, and there are only few teams in the world who are tackling this problem. The official release of Super Zero (BERG) to the world presents the first anonymous cryptocurrency that supports smart contracts. The BERG's R&D team (BERG Team) is the only team in the world that presents a complete solution to solve the Privacy problem and has completed major R&D work. BERG team not only considers the privacy of DAPP Users' accounts and transactions but also fully considers privacy protection of DAPPs' developers, making privacy protection of the DAPP Ecosystem truly secure and stable.



BERG team has assembled a 3 in 1 suite that can provide a complete privacy protection solution for DAPPs; including advanced innovative technology components BERG (privacy cryptocurrencies platform supporting smart contracts), ALIEN protocol (a protocol that can solve security problems within the transmission of information in decentralized networks) and CASTROL protocol (a protocol that protects decentralized networks and provides privacy protection for every node in the Internet). The white paper describes BERG's work and includes core information about the project as well as the disclosure of subsequent project plans.



## CHAPTER II INTRODUCTION

### 2.1 DECENTRALIZATION TECH AND PRIVACY ISSUES

At present, users have an increasing concern and demand for privacy protection, many well-known companies have leaked a large number of user privacy data, including Yahoo, Uber, PayPal, InterContinental Hotels Group, US credit agency Equifax, UK National Health Service System(NHS) etc., compromising tens of millions to hundreds of millions of user data. Facebook lost tens of billions of dollars in market value in two days due to one of the largest privacy leaks in March 2018. The issue of privacy has also attracted the attention of many governments, the European Union took the lead in promulgating the General Data Protection Regulations (GDPR) to urge companies to effectively protect users' privacy.

Majority of the privacy leaks in the Internet application scenarios are caused by the lack of adequate data security protection mechanisms in a centralized platform. Blockchain technology is thought to be able to prevent such incidents. The design of blockchain networks such as Bitcoin and Ethereum didn't take into account the possibility of the link established between the wallet and physical identity. The extremely sensitive information such as digital assets and their transaction records in the blockchain is transparent to public and cannot be tampered with. If blockchain is used in a larger number of real scenarios, the transparency is undoubtedly unacceptable for most users.

The range of legal use cases of financial privacy is very wide. Financial privacy protection is needed for most transactions in the world. It is unreasonable to expose cryptocurrencies' assets and transactions data stored on the blockchain to the public.

#### Examples of real-world scenarios

- \* A company wants to protect supply chain information without revealing it to the competitors
- \* An individual does not want the public knowledge of paying for consultation with a bankruptcy lawyer or divorce lawyer
- \* A family, fearing discrimination, wants to withhold children's medical history from employers and colleagues

\*A wealthy individual preventing potential criminals from gaining access to his whereabouts to prevent extortion.

\* Commodity buyers and sellers want to avoid the transaction being cut off by any middlemen.

\* Investment banks, hedge funds and other types of entities dealing with trading financial instruments (securities, bonds, derivatives); protecting their positions or trading intentions.

In smart contracts, the entire sequence of actions is distributed through the network and recorded on the blockchain and is publicly visible. Individuals and organizations believe financial transactions (such as insurance contracts or assets transactions) are highly confidential; however, this need for the information privacy protection is not currently supported. The lack of privacy becomes the main obstacle to the widespread adoption of decentralized smart contracts. The lack of privacy protection technology is a serious bottleneck for the popularization of DAPPs. The technological development progress in related fields has attracted public attention.

## **2.2 BLOCKCHAIN PRIVACY RISK**

Bitcoin network is a typical blockchain technology representative. Mainstream cryptocurrencies in the market are mostly based on the same technical features. The following uses Bitcoin network as an example to analyze the risk of privacy leakage.

First, from the perspective of bitcoin transaction system's structural design:

\* UTXO model of transaction data contains input address and output address information, each input address points to the previous transaction, and all input transaction amount can be traced back to the source.

\* Transaction data is stored in a public global ledger, and any participating user can obtain a complete global ledger. In the consensus process, the verification node needs to retrieve historical transactions, and all transaction information is not encrypted to protect data.

The addresses of participants in bitcoin transactions are created by the users and not related to the identity information. No one can directly deduce the identity information of the users in the transaction by observing the transaction records; however, there is a correlation between transactions disclosed in the global ledger, and potential attackers can deduce the transaction rules of bitcoin addresses by analyzing the transaction records in the global ledger. The transaction frequency, transaction characteristics, and correlation between addresses, etc. in the global ledger makes it possible for an attacker to associate a bitcoin address with the identity of a particular user in the real world.

One of the methods mainly obtains regular characteristics of the transaction of the address by analyzing the transaction record related to the address and estimates the identity information of the corresponding user accordingly. Since there are unique transaction characteristics in certain types of blockchain transactions, attackers can restore the actual transaction according to the transaction characteristics of the address, thereby determining possible identities of the user. Androuraki E. et al

designed a simulation experiment to match the blockchain address with the students' identities. Students used bitcoin as a payment method for daily transactions and used the one-time address method recommended by bitcoin to enhance privacy protection. Analysts were able to successfully match the student's identity and the blockchain address with 42% accuracy through behavior-based clustering technology. Monoco J. V. et al quantifies the trading behavior of bitcoin users and analyzes the trading rules of users based on twelve dimensions including trading time, interval and cash flow. After 6 months of experiments, the accuracy of using the analysis model to successfully identify users' real identities is as high as 62%, and the error rate is less than 10.1%.

Another method is to use some potential knowledge in blockchain transaction designed to cluster different addresses and get multiple addresses of the same user.

Currently, there are mainly three rules for address clustering:

\* For a transaction with multiple input addresses, it is generally assumed that all input addresses come from the same user or a collection of users. When a user initiates a transaction, the digital assets may come from multiple addresses of the user, and the user needs to sign each input address separately, so most multi-input transactions come from the same user. The rule has been applied to many research projects and has achieved good clustering results.

\* In the transactions organized by the mining pool, multiple output addresses in the same transaction belong to the same user group. As the difficulty of "mining" increases, individual "miners" are no longer able to win the competition, requiring hundreds of "miners" to join the "mining pool" to complete "mining" together, and the rewards will be distributed to the "miners" participating in the collective "mining".

\* The change address and the input address in the transaction belongs to the same user. In one transaction, the total amount in the input address may be larger than the amount issued by the user, so the bitcoin system will automatically generate a change address for the sender to receive the change amount in the transaction. As with other addresses, the change address may be selected by the system as the input address in the new transaction, the output address will usually only be used once. Since the change address was regenerated by the system during the transaction, it is impossible for an address to be used as the input address and output address for one transaction at the same time; there must be another output address other than the change address in the output of the transaction. By using the characteristics of change addresses, we can figure out the relationship between other addresses.

Studies have found the relationship between many addresses in bitcoin system using the above clustering rules. Meikle John S. et al attained the identification of bitcoin addresses in bitcoin theft cases by using heuristic clustering methods. Dmitry E. et al also provided a method to automatically identify cluster bitcoin addresses.

### **2.3 DECENTRALIZED PRIVACY PROTECTION TECH**

We are pleased to see teams are beginning to address privacy protection of decentralized networks; including Zcash, Monero and Dash. One widely used method is to change the transaction process without changing the transaction results, and attackers cannot directly obtain complete

information about the transaction. The method is called " mixed currency". For example, in Chaum D.'s article, an anonymous communication technology is mentioned, which hides the real communication content in the communication process. The basic idea can be expressed by formula (1):

$$CM(Z1,CA(Z0,m),A) \rightarrow CA(Z0,m),A(1)$$

The left side of equation (1) is the message sent by the sender to the intermediary, and the right side is the message sent to the receiver after the information is processed by the intermediary. The sender wants to send the messages  $Z0$  and  $m$  to the address  $A$  of the receiver. First, the message encrypts with the key  $CA$  of the receiver to obtain  $CA(Z0, m)$ , then packages the authentication message  $Z1$  of the intermediary. The encrypted message  $CA(Z0, m)$  and the address  $A$  of the receiver, then encrypts with the public key  $CM$  of the intermediary to prevent the information from being intercepted or tampered with by attackers during the sending process. After receiving the information, the intermediary decrypts it with his private key to get  $Z1, CA(Z0, m), A$ , but is unable to decrypt the content of  $CA(Z0, m)$ . The intermediary sends  $CA(Z0, m)$  to address  $A$  after verifying  $Z1$  is correct. The receiver then decrypts the message using its own private key to complete the communication.

Messages are not directly transmitted between the sender and the receiver, instead, the messages are transmitted indirectly through an intermediary, making it impossible for attackers to observe the communication behavior between the sender and receiver, thus, improving the anonymity of the communication. If the message is passed through multiple intermediaries, for the difficulty for attackers to discover the communication relationship between the sender and receiver increases.

The mixed currency mechanism in cryptocurrencies draws from the above methods ( Dash and Monero ) and removes the traceable relationship between the actual sender and receiver in the transaction through an intermediate hierarchy. The implementation of the currency mixing process can be implemented by a trusted third-party or other protocol. A third-party node is involved in the currency mixing process, the existing currency mixing mechanisms can be divided into two categories: the central node and the decentralized node. The two mechanisms have their own advantages and disadvantages in terms of currency mixing reliability, efficiency and cost.

More sophisticated encryption techniques have been applied to blockchain privacy protection with the development of technology - Zcash using Zero-Knowledge proof.

Brief descriptions of the three most popular cryptocurrencies with privacy protection:

### **Zcash**

Zcash is a cryptocurrency that use encryption technology to provide users with greater privacy protection than other cryptocurrencies such as Bitcoin. Originally named ZeroCoin, the team then developed the ZeroCash system, and it was developed into Zcash cryptocurrency in 2016.

Zcash payments are published on the blockchain, users can use optional privacy function to hide the sender, receiver and amount of transactions on the blockchain. Only those who have the key can see the contents of the transaction. The user has full control and can choose to provide the key to others to prove payment for auditing purposes.

Zcash is an improvement to Bitcoin and developed on Bitcoin's infrastructure. It uses Zero-Knowledge Proof technology called zk-SNARKs to encrypt user information. Zk-SNARKs is an encryption method based on pure mathematical theory. The encryption method is self-contained, it has the advantage of not depending on external operating environment; therefore, has a wider range of application scenarios.

Since Zcash uses the same underlying architecture as Bitcoin's network, it can only support simple transactions, similar to the Bitcoin network with a pre-set privacy protection mechanism. Using Zero knowledge for the encryption of transactions is inefficient, and the application scenario is restricted further.

### **Monero (XMR)**

Monero was founded in April 2014. Unlike Zcash, Monero did not choose to develop a blockchain system based on Bitcoin. The design is modular from the bottom and has good scalability.

Monero features the "proof of work" (POW) consensus mechanism. Unlike many previous cryptocurrencies, the Monero's algorithm CryptoNight is an AES intensive and memory-consuming operation; which significantly reduces GPU's advantage over CPU - reduces the risk of centralization of POW.

Monero uses Ring Confidential Transactions algorithm for its encryption method. The method mixes the signer's public key with another public key set and then signs the message, makes it impossible for intruders to distinguish which public key corresponds to the actual signer; therefore, protecting the user's real identity. Monero's mixed-currency participating users do not need to communicate with other participating nodes, they can participate in mixed-currency by themselves, providing effective protection measures for common distributed denial-of-service (DDOS) attacks and information disclosure in the decentralized mixed-currency mechanism.

Monero does not support smart contracts, and the high risk of being attacked is still present even though it adopts decentralized mixed currency technology. Users need to rely on the public key of other users when using Ring Confidential Transactions technology. If other users are malicious, the problem of users' privacy disclosure will arise.

### **Dash (DAS)**

Dash is the first cryptocurrency designed to protect privacy. Dash utilizes centralized currency mixing scheme – simply transfer a sum of money to multiple addresses several times; simple to implement and easy to operate; no other technical support is needed during its currency mixing process. The centralized currency mixing scheme has high applicability in various cryptocurrency systems; however, the existing solution requires sender and receiver to mix coins online. If the sender and receiver can't reach an agreement on the amount of mixed currency, the transaction must be postponed. There is a common delay problem and the currency mixer is centrally deployed. The nodes of currency mixer can obtain all the information of the transaction and able to pilfer cryptocurrency. The most improved schemes prevent theft and information disclosure by increasing the cost of third-party violation; fundamentally, cannot eliminate the occurrence of violations. The mixed currency



scheme using cryptography techniques such as blind signature increases the calculation cost. The execution of the mixed currency process by a third-party will inevitably bring additional service costs.

Dash does not support smart contracts, and the third-party mixed-currency provider mechanism relies on the credibility of the third-party and encounters unpredictable risks. In recent years, Dash has focused on the development and layout of ecological applications based on its good circulation and has strengthened cooperation with enterprises, attempting to make Dash a payment tool with strong circulation value instead of the emphasis on privacy protection.

### **Conclusion**

From the aspect of the latest technology, solution to ensuring privacy protection by the adaptation of the latest cryptographic algorithms - the non-interactive zero-knowledge proof mechanism (NIZK), shows the most promising improvement. The use of encryption mechanism requires significant changes to the underlying protocol and consumes more computing resources, affecting the efficiency of blockchain applications. Therefore, the usage of privacy protection mechanism needs to consider the efficiency and cost of nodes in efficiency performance, and cost of computing and storage.

In the decentralized application scheme, smart contracts widely increase the application scenario of blockchain. The applications are no longer limited to the digital assets value of circulation. The current mainstream blockchain privacy protection technology does not support smart contracts, which prevents the greater establishment of practical usages. Any secure privacy protection mechanism for anonymity to support smart contracts must make major modifications to the underlying system of blockchain, the implementation will be difficult.

BERG is the solution to solve the above problems.

## **2.4 OVERVIEW OF BERG SOLUTIONS**

BERG (Super Zero) is the world's first blockchain system that truly realizes the complete privacy protection of blockchains through non-interactive zero-knowledge proof. Compared to the existing blockchain privacy protection technologies, BERG not only can realize the privacy protection of account and transaction information but also support Turing complete smart contracts. In addition, developers can also create their own encrypted cryptocurrencies supporting smart contracts based on BERG-Chain.

BERG re-designed the blockchain structure and various underlying protocols, making Turing complete smart contract for privacy protection come true. Making privacy protection measures available for a wider range of application scenarios, and making the attacks on user's private data more challenging with the advanced NIZK encryption algorithm. In addition, the upcoming BERG V1.0 release, NIZK encryption algorithm is thoroughly optimized, which greatly reduces the memory resources required and improves the computational efficiency. Compared with the mainstream privacy cryptocurrencies, BERG's supports of Turing complete smart contracts, privacy protection measures and its related decentralized applications have significantly broadened its use-case scenarios.

More importantly, the BERG team considers the privacy protection measures required by the decentralized applications. The team also plans to provide solutions for the security of point-to-point network transmission and the privacy of the physical network address of the account, enables the centralized application to obtain powerful privacy protection functions when interacting with the centralized application or when interacting with the user's client.

The entire integrated solution will consist of a complete set of 3 in 1 suite, where BERG is the first publicly released project and the other two projects positioned as following:

ALIEN Protocol: A distributed DNS system that can use existing P2P network interaction information, has the functions of IP automatic switching and dynamic addressing, resists attacker blocking, and enables the entire data transmission network to achieve the ideal stable security.

CASTROL Protocol: The anonymous protection of IP addresses can be realized through decentralized network, which can be used to protect the privacy of physical nodes in both centralized and decentralized networks.





BERG

## CHAPTER III THE DESIGN OF BERG

### 3.1 DESIGN PRINCIPLES

The privacy protection technologies of decentralized network in the existing market do not combine with decentralized applications, particularly, the implementation of smart contracts is not protected. The sequence of actions performed in the smart contract is publicly visible throughout the network and / or recorded on the blockchain platform. In Turing complete blockchain network, BERG's design must meet several basic principles as well as meet the system's capacity requirements.

**Un-traceable** - every transaction in the blockchain network has an input and an output, constructing an acyclic graph of transactions, on which all of the transaction flows can be tracked, all of the transaction sequences can be concatenated and traced. BERG is designed to break the link between the two transactions, making the attack impossible.

**Un-associable** - each user in the blockchain network has their own collection address. Once the address is associated with the real user identity, all the transactions occurring at the address in the network can be associated with the corresponding user identity, resulting in the exposure of the associated behaviors to the address. All the transactions and balances are still publicly visible when a user creates a new pseudonym public key for anonymity. BERG uses encryption technology to make the payment address unrelated.

**Anti-statistical analysis** - actual user behavior has statistical characteristics. If the transaction data in the blockchain network has a correlation that reflect such statistical characteristics, it is possible to deduce the addresses belongs to a specific user through statistical analysis of the blockchain data. When ring signatures are used, the ability to resist statistical analysis will decrease if ring members or nodes are malicious. BERG must be able to completely hide the address and the relationship between addresses by technological means.

**Practicality principles** - BERG, while hiding the transaction data, will not take all the information into its scope, which can be uneconomical and inefficient. BERG will consider the user's existing usage habits and concerns to carry out research and development periodically.

**Optional auditing solution** - for the alternative audit scheme and certain complex business applications, the user may choose a trusted third-party to conduct financial audit of transactions. The user should have the ability to give the third-party to track the specific information from the transactions.

### 3.2 IMPLEMENTATION PLAN

In the first phase, BERG will completely protect the inputs and outputs of the trading system and the trading details through non-interactive zero-knowledge proof (NIZK). The transaction details are invisible to everyone except the two parties involved. BERG will maintain the smart contracts running on the chain and integrate the assets generated by the smart contracts with BERG's own trading system, considering that the online running smart contracts and the total number of open contracts issued assets have universal applicability. This will enable the privacy of the assets generated by the smart contract.

In the second phase, within the smart contracts running online, BERG will provide a latent structure called Hidden Data Structure(HDS) to satisfy the requirement for the total number of issued assets with protected contracts. The calculations for the HDS complete off the chain. The function will protect the total number of contractually issued assets.

In the third phase, BERG will adopt a more advanced consensus mechanism to improve the throughput of BERG networks. At the same time, BERG will decompose the operation of the contract into two steps: offline calculation and online verification. The offline calculation will fully understand the calculation rules and data, and will return the encrypted result. When the result is submitted online, the online node will only validate the result and determine whether the data conforms to the calculation rules; the node will not know the details of the data and calculation rules.

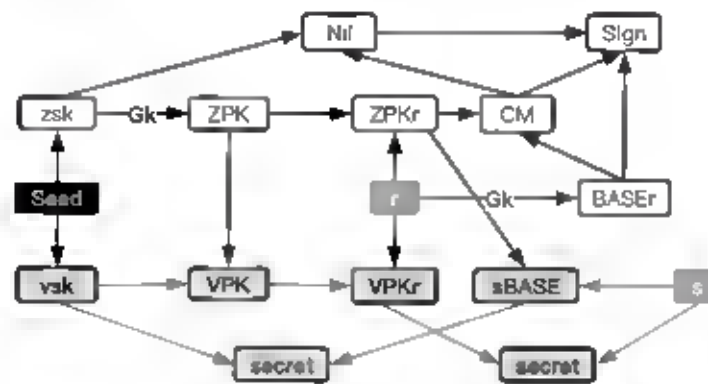
### 3.3 BERG PROTOCOL

#### Account System

Accounts are divided into two categories: user account and contract account. The user account is a 32-byte *seed* selected by the user, the contract account generates a 64-byte *address* corresponding to the smart contract environment the user installed; both categories are unique and non-repeatable.

The user account can generate a 64-byte private key *SK* and a 64-byte public key *PK*, as the user's payment address. When installing or invoking the smart contract, the wallet will generate a temporary address  $PK_r$  according to the current condition. The temporary address cannot be associated with the user's private key and public key and will only be used once.

When the smart contract is installed, the wallet will change the temporary address to a 64 byte smart contract address (*CADDR*) in accordance with the current condition. As the node receives the address, it needs to ensure that the contract address has not appeared before.



Let

$G_k = \text{NewEcc}()$

$seed = \text{New}(\text{Byte32})$

$r = \text{RandFr}()$

$s = \text{RandFr}()$

$n = \text{RandFr}()$

$m = \text{Message}()$

SK

$zsk = \text{HASH}_{zsk}(seed)$

$vsk = \text{HASH}_{vsk}(seed)$

$sk = (vsk \ zsk)$

$zvsk = zsk \ vsk$

PK/TK

$ZPK = zsk \ G_k$

$VPK = vsk \ zsk \ G_k$

$PK = (ZPK \ VPK)$

$TK = (ZPK, vsk)$

PK,

$$\begin{aligned}
 BASE_r &= r \cdot G_k \\
 ZPK_r &= r \cdot ZPK \\
 VPK_r &= r \cdot VPK \\
 PK_r &= (VPK_r, ZPK_r, BASE_r)
 \end{aligned}$$

*Trace*

$$VPK_r = vsk \cdot ZPK_r$$

*Enc*

$$\begin{aligned}
 BASE_s &= s \cdot ZPK_r \\
 SECRET &= s \cdot VPK_r \\
 key &= Hash_2(SECRET) \\
 M &= Enc_{sk}(m, key)
 \end{aligned}$$

*Dec*

$$\begin{aligned}
 SECRET &= vsk \cdot BASE_s \\
 m &= Dec_{sk}(M, key)
 \end{aligned}$$

*Sign*

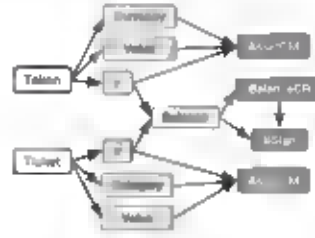
$$\begin{aligned}
 k &= Hash_1(a, zvsk \cdot m) \\
 s0 &= k \cdot BASE_r \\
 h &= Hash_2(S0, m) \\
 s1 &= k + zvsk \cdot h \\
 sign &= (s0, s1)
 \end{aligned}$$

*Verify*

$$s1 \cdot BASE_r = S0 + h \cdot VPK_r$$

*seed* is an account seed and users must keep it securely. *sk* is a private key and cannot be stored persistently. *TK* is a tracking private key and can be provided to a trusted third-party for account auditing. *PK* is the public key and provides the transaction destination address to other users. *PK<sub>r</sub>* is a temporary address, provided to the smart contract, and temporarily used to receive the asset.

**Assets System**



Let

$$G_r = \text{NewEcc}()$$

$$H_{for\_s} = \text{Hash}_{for\_s}()$$

Token

$$token = (value, Currency)$$

$$G_{cy} = \text{Ecc}(Currency)$$

$$r_{tkn} = \text{RandFr}()$$

$$CM_{asset,tkn} = value \cdot G_{cy} + r_{tkn} \cdot G_r$$

Ticket

$$ticket = (value, Category)$$

$$G_{cy} = \text{Ecc}(Currency)$$

$$r_{tkn} = \text{RandFr}()$$

$$CM_{asset,tkn} = value \cdot G_{cy} + r_{tkn} \cdot G_r$$

Balance

$$r_{balance} = \sum_{i=0}^n r_{in} - \sum_{i=0}^m r_{out}$$

$$CB = r_{balance} \cdot G_r$$

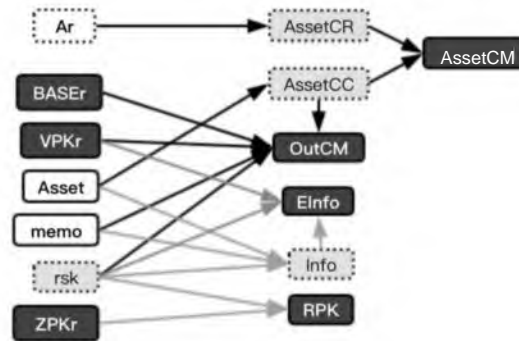
$$sign_{balance} = \text{Sign}(G_r, H_{for\_s}(r_{balance}))$$

$$check_{sign} = \text{Verify}(G_r, CB, H_{for\_s}(sign_{balance}))$$

$$check_{balance} = \sum_{i=0}^n CM_{asset,in} - \sum_{i=0}^m CM_{asset,out} = CM_{asset,balance} + CB$$

User account or smart contract account, has the attribute of managing unlimited variety of assets. With the exception of the settlement of transaction fees using BERG coins, each asset has the same transaction characteristics as BERG coins. Excluding BERG coins, the remaining assets can be generated by a smart contract. Each asset can be given a name of up to 32 byte length (token name) for mnemonic purposes and these names are also not allow to ben reused. The asset type can be specified when the account performs balance queries or transfer operations.

### Output Construct



$$ar = RandFrQ$$

$$rsk = RandFr()$$

$$C\&sser = ar \cdot Gcr$$

$$CCasset = asset \cdot Gcc$$

$$CMasset = CRasset + CCasset$$

$$CM_{CM} = Hashp\{CCasset, BASEr, VPKr, memo, rsk\}$$

$$Info = Append^{asset, memo, rsk}$$

$$EInfo = Enc^{nfo, rsk}$$

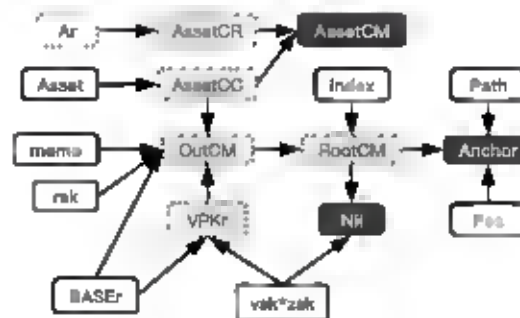
$$RPK = rsk \cdot VPKr$$

$$Inputs = (CMasset, CM_{out}, EInfo, PKr, RPK)$$

$$Vars = (ar, asset, memo, ZPK, rsk, CR, CC)$$

$$Proof = Prove^{Inputs, Vars, CIRCUIT()}(output)$$

### Input construct



$$CM_{out} = Hash_{pedersen}(index \cdot CM_{in})$$

$$Anchor = MerkleRoot(index \cdot CM_{out} \cdot Path)$$

$$Nil = z_{vsk} \cdot CM_{out}$$

$$Tid = vsk \cdot CM_{out}$$

$$Inputs = (CM_{in}, Anchor, Nil)$$

$$Vars = ([the \text{ resi}])$$

$$Proof = Prove(Inputs, Vars, CIRCUIT_{input})$$

$CM_{out}$  is the value of the leaf node of the Merkle tree composed of UTXO sequences.  $Anchor$  is the root of the Merkle Tree where the input data is located.  $Path$  and  $Pos$  are certification paths from  $CM_{in}$  to  $Anchor$ .  $Nil$  is the hash of 32 Byte for destroy the used  $OUT$  of UTXO.  $Tid$  is the hash of 32Byte for transaction trace.  $CM_{out}$  is the output commitment of the transaction.  $CM_{in}$  is the asset commitment of the input.

### License System

#### LICENSE

$$Inputs = (ZPK, prop)$$

$$sk_{license} = (sk_{zpk}, sk_{prop})$$

$$PK_{lic} = (PK_{zpk}, PK_{prop}) = (sk_{zpk} \cdot G_{lic}, sk_{prop} \cdot G_{lic})$$

$$r_{lic} = Rand(Fr)$$

$$R_{lic} = r_{lic} \cdot G_{lic}$$

$$S_{lic} = r_{lic} + sk_{zpk} * Hash(ZPK) + sk_{prop} * prop$$

$$LIC = (prop, R_{lic}, S_{lic})$$

**PROVE**

$$R + Hash(ZPK) \cdot PK_{zpk} + prop \cdot PK_{prop} = s_{lic} \cdot G_{lic}$$

In BERG's Alpha and Beta networks, in order to ensure the healthy development of the network at the initial stage, to ensure the robustness of the consensus and the timeliness of system updates, it is necessary for BERG project team to coordinate the miners. Therefore, testers with mining needs need to apply for mining licenses from the BERG R&D team. In addition to mining, testing of other functions does not require a license. On the premise of not disclosing the identity of miners as much as possible, the block will expose some of the attributes in the license, which can be monitored by the BERG community. In the early stage of the Beta network, when the network is attacked and a major crisis occurred, the BERG team will use unconventional means to resist attacks and protect the property safety of community members through community voting under the premise of community permission and supervision. The license feature will be removed after BetaNet has been online for half a year

**Witness System**

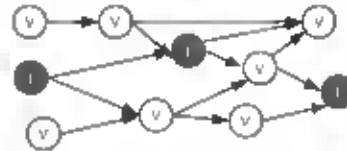
BERG protocol uses non-interactive zero knowledge proof (NIZK) and needs to provide witness information of assets source when generating transactions. Each node will verify according to the witness information. BERG uses the Merkle Tree to maintain a witness system that records status changes. The system will provide verification function at the nodes and authentication information at the wallet side

$$ROOT = MerkleRoot(POSITION, LEAF, PATH)$$

*ROOT* is the root of the current Merkle Tree. *LEAF* is the leaf at this *POSITION*. *PATH* is the proof path from *LEAF* to *ROOT*.

**Prove System**

BERG's proof system includes a calculation circuit based on directed acyclic graph to describe the internal constraints of each BERG transaction: input and output balance of various asset types, verification of public and private key, validity of commitments, validity of witness, etc. The circuit loaded with data can generate a *Proof* through non-interactive zero knowledge proof (NIZK). From the submission of the *Proof*, nodes can verify various parameters and constraints loaded in the circuit while hiding a large amount of detailed information





$$\begin{aligned} Inputs &= (I_0, I_1, \dots, I_n) \\ Vars &= (V_0, V_1, \dots, V_m) \\ Proof &= Prove(Inputs, Vars, CIRCUIT) \\ Check &= Verify(Inputs, Proof, CIRCUIT) \end{aligned}$$

$I_i$  is the public data in the transaction.  $V_j$  is the privacy data in the transaction. All variables construct a *CIRCUIT*. Prove system generate the *Proof* with *Inputs*, *Vars* and *CIRCUIT*. Then carry *Inputs*, *Proof* in the transaction. the verification process is verified by *Inputs* and *CIRCUIT*, *Proof*.

#### Process Step

##### A. Compute

The user calculate *RESULT* uses the information provided by the account and witness system, and input data according to the current required calculations. The calculation rules runs under off chain environment to obtain the results.

$$RESULT = COMPUTE(METHOD, ACCOUNT, DATA, WITNESS)$$

##### B. Prove

The user generates *STX* with the random variable  $r$  and the result of compute, and submits it to the node. The *STX* include check data  $C_i$ , encode-data  $E_j$  and proof-data  $P$ .

$$\begin{aligned} STX &= PROVE(RESULT, r) \\ STX &= ((C_0, C_1, \dots, C_n), (E_0, E_1, \dots, E_m), (P_1, P_2, \dots, P_m)) \end{aligned}$$

##### C. Verify

After receiving the *STX*, the node confirms  $C_i$  in the witness system and the proof system. When certificates are verified as correct, the node will accept the *STX*.

$$\begin{aligned} ret_i &= VERIFY(C_i) \\ Check &= ret_0 \& ret_1 \& \dots \& ret_n \end{aligned}$$

##### D. Confirm

As the asset receiver synchronizes to the verified transaction *STX*, the receiver uses the private key to decipher the encode-data  $E$  and generate plaintext  $D_i$ . The plaintext  $D_i$  and proof  $P_i$  are input to the proof system for verification, success indicates the transaction is validated. When the transaction is confirmed by  $n$  blocks, the transaction recipient considers the transaction has been confirmed.

$$D_i = FETCH(E_i, ACCOUNT)$$

$reti = CONFIRM^{\wedge}$  (修)

$Check = ret_e \& ret_x \dots \& ret_m$

The execution steps of BERG are open, the abstract description of the steps and parameters support the new functions from phases one to three described in the "Implementation Plan" section with minimal changes to the code structure during subsequent upgrades.

#### **General Privacy Transaction**

Within BERG, data in ordinary transactions are encrypted; non-trading parties will not know the details of source, destination, asset type, amount, etc. The system does not distinguish between assets generated by smart contracts and BERG's own assets in transaction processing.

#### **Online Smart Contract**

BERG's General smart contract can be used for public calculations, to develop statistical plans, disposal rules and publicity rules for various assets; the input and output information must be isolated from users' real identities using temporary addresses.

BERG smart contracts are compatible with Ethereum smart contract instructions; most of Ethereum smart contracts can be run on BERG without modification.

#### **Online Privacy Assets**

Smart contracts issue assets by calling the online privacy asset issuance methods. The number of issued assets is visible to all users and has transaction attributes equivalent to BERG's own assets, which can be handled through private transactions.

#### **Offline Privacy Assets**

Smart contracts issue assets by calling the offline privacy asset issuance methods. The number of issued assets is not visible to users. The assets have transaction attributes equivalent to BERG's own assets and can be processed through private transactions.

#### **Offline Smart Contract**

BERG's offline smart contracts run only on the user's machine. The calculation rules are only visible to some users, and the consensus verifies the correctness of the results.

### **3.4 ABOUT NON-INTERACTIVE ZERO-KNOWLEDGE PROOF (NIZK) PERFORMANCE OPTIMIZATION**

For the blockchain system adopting the non-interactive zero knowledge proofs ( NIZK ) scheme, the biggest application bottleneck at present is that the time for generating proof at transaction is too

long. The zero knowledge proofs module Super-ZK of BERG system have made the following breakthrough innovation for this bottleneck.

1. Currently, we use the zk-SNARKs framework to generate NIZK, which uses the ALTBN128 curve and the Groth16 preprocessing process, which reduces the computation time by 1/3 compared with PGHR13 preprocessing scheme. Although the zk-SNARKs framework requires a credit installation process, computing circuits will not be dynamically constructed in the BERG implementation, so the zk-SNARKs framework can meet the BERG requirements in all current scenarios.

2. We have innovatively developed a Twisted Edwards curve to replace the SEIA256 to generate the public key, and ECC Hash to generate Merkle tree, which can increase the transaction speed by more than 4 times.

3. BERG adopts a single input and output structure, and each description is linked by an asset channel. Such a circuit configuration is more modularity, in the case of multi-core CPU, the efficiency of parallel execution is increased three times.

4. Part of the codes of Super-ZK is written by assembly language, which optimizes engineering structures such as resource allocation and makes it more efficient in code execution.

Based on the above optimization, compared with other blockchain systems that directly use zk-SNARKs, we have an order of magnitude speed increase in transaction proof generation speed, which greatly improves the applicability of the NIZK system.

### **3.5 THE PRINCIPLE OF SMART CONTRACT ISSUANCE AND OPERATION OF ANONYMOUS ASSETS**

#### **The Models of UTXO and ACCOUNT Model**

The blockchain is a distributed ledger, the smallest unit of the ledger account is record, and each record stores the inflow or outflow of the account's assets. According to the different ways of asset outflow recording, the blockchain systems have evolved into two different accounting implementations, which are UTXO model and ACCOUNT model. These two modes correspond to the models of Bitcoin and Ethereum respectively.

The advantages of UTXO model is:

A. Each transaction in UTXO model is independent of each other, which means that transactions under one account can be processed in parallel as long as the problem of double spend can be handled properly, and the capability of multi-core CPU can be fully utilized.

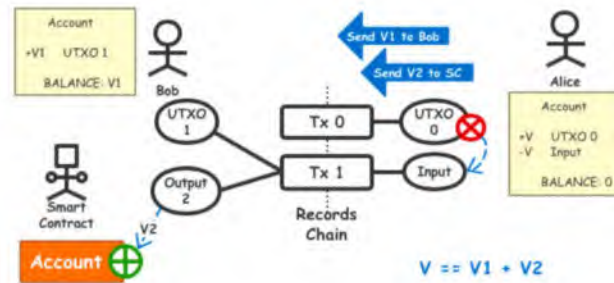
B. UTXO is essentially a record form based on history, both a process and a result, so it has great advantages in some applications where proofs need to be generated. This is why the blockchain systems with CT(Confidential Transactions) features are basically UTXO mode.

The advantages of ACCOUNT model is:

A. ACCOUNT model directly add or subtract the assets in an independent account, and can add or subtract any value of assets in an account with only one record. Therefore, the generated record size is much smaller than the record generated by UTXO in the same situation.

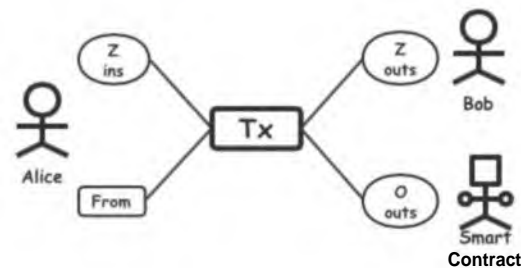
B. The ACCOUNT model is state-based, *Input* and *Output* are processes, and *ACCOUNT* is the result, so it is easy to implement the Turing machine, that is why the blockchain that supports Turing's complete smart contracts mostly uses the ACCOUNT model.

### BERG's Hybrid Model



BERG applies UTXO and ACCOUNT models together, using UTXO model where CT(Confidential Transactions) are required, and using ACCOUNT model where smart contracts need to be run. BERG seamlessly integrates these two models through trading, consensus, and Pedersen Commitment algorithms, enabling smart contracts to perform surprisingly.

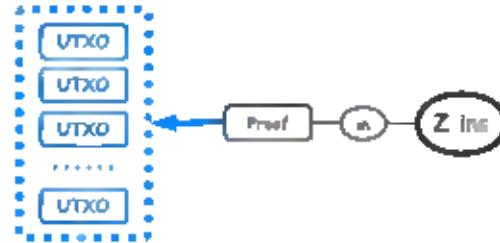
### Structure Supporting Anonymous Transactions with Smart Contracts



#### A. Transactions $Tx$

BERG's confidential transactions  $Tx$  have an anonymous input set  $Zins$ , an anonymous output set  $Zouts$ , a normal output set  $Oouts$ , and a temporary address  $From$ .  $Zins$  is completely anonymous, making it impossible for third-party observers to find out the source and content.  $Zouts$  is completely

anonymous UTXO", only the receiver can view and use its content, the content carried by the  $O_{out}$  is not hidden, and it refers to the receiver in two situations: one is pointing to a smart contract address, and the other is pointing to a temporary address.  $From$  represents the sender of the transaction and is also a temporary address. Therefore, the whole  $Tx$  cannot make the third-party observers determine who is the real user, and the information such as assets carried in it is also hidden to the greatest extent.



#### B. Input $Z_{in}$

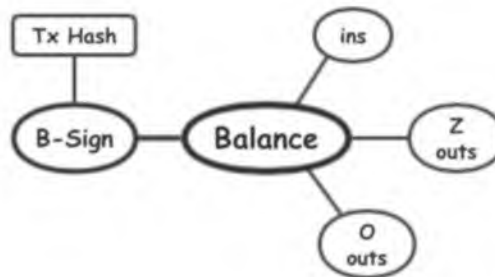
In the input set  $Z_{in}$  of the BERO transactions, each input is anonymous, including the ID of the source "UTXO" and the asset information carried. Each input is converted to a  $Proof$  generated by zero-knowledge proofs  $ZKP$ . The  $Proof$  points to a particular UTXO hidden in a huge UTXO sequence, which is part of the BERO history, all details hidden by  $Proof$ . The verifier can confirm whether the input is valid by  $Proof$  without knowing the details.

#### C. $Z_{out}$ Output and $O_{out}$ Output



$Z_{out}$  points to the temporary address  $PK_r$ , and the temporary address can only be decrypted by the receiver. Since each temporary address is different, no third party can identify the  $Z_{out}$  real point address.  $Z_{out}$  also carries the encrypted information  $Encrypt Info$  of the asset, which can only be decrypted by the party holding the receiver's private key. And  $OutCM$  is the output commitment; only the parties of the transaction can reproduce the computing process of  $OutCM$ .  $OutCM$  plays a crucial role in proving that  $Z_{out}$  is referenced by  $in$ .

The  $PK_r$  pointed to by  $O_{out}$  has two forms: one is initiated by a smart contract and points to the temporary address of the common account. The other is initiated by a normal account, pointing to the address of the smart contract. Due to the randomness of the temporary address, the third party cannot know the identity of the receiver.



D. The *Balance* of inputs and outputs

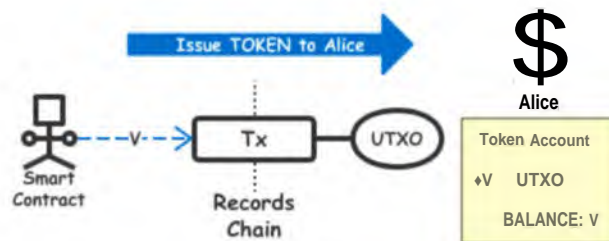
*Tx* packages *ins*, *Zouts*, and *Oouts* together. How to prevent malicious attackers from tampering with the data and ensuring asset security? By introducing Pedersen Commitment, its homomorphic encryption feature enables verifiers to confirm that *Balance* must be balanced without knowing the details of the information, that is, the input is equal to the output.

In addition, in order to prevent malicious attackers from tampering with *Oouts*, we use the random feature of Pedersen Commitment to sign *Tx Hash* with the random part of *Balance*. In this way, each input and output can be computed independently and then packaged together by *B – Sign*.

E. Transaction Sender *From*

When the output of a transaction points to a smart contract, the smart contract sometimes needs to output resources to a given account based on the written rules. At this time, the temporary address *From* is the place where the output resources are received. *From* is determined when the transaction is generated and used only once. Other than the sender of the transaction, the third parties cannot locate the identity of the sender.

F. Issuance of Anonymous Tokens

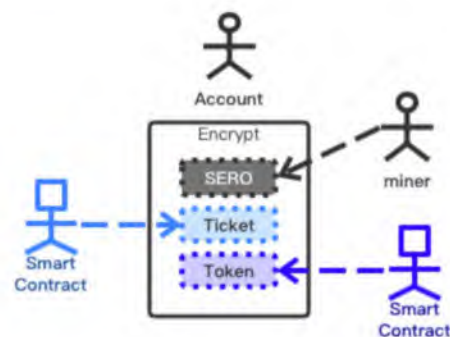


Token, also known as the “homogeneous pass”, is an asset form recognized within BERG system.

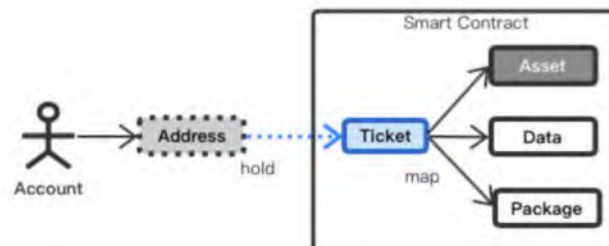
The same kind of tokens can be arbitrarily divided and mixed. Tokens on BERG system is different

from Ethereum's tokens. As the first currency of the BERG system, BERG coins are also a type of tokens in essence. For Token assets, except for handling fees, which can only be paid in BERG coins, they are treated the same within the BERG system, and their privacy and security are ensured by consensus. Once the anonymous Tokens are successfully issued, the smart contract can send the tokens to the temporary address  $PK_T$  of a normal account in the form of a normal transaction. At this time, these sent tokens will be separated from the smart contract account in the form of UTXO, and transfer into the user's account like BERG coins, thus being protected by BERG's privacy mechanism.

#### G. Issuance of Anonymous Tickets



Tickets, also known as the “homogeneous pass”, is another more extensive asset form recognized within BERG system. Different from tokens, tickets are an inseparable form of assets that is unique to the individual. Like tokens, BERG also provides anonymity support for Tickets. When anonymous tickets are issued, the tickets are represented as an ID that are sent as the normal transaction to an account's temporary address  $PK_T$  and the tickets are pulled separated from the smart contract account as UTXO and protected by BERG's zero-knowledge proofs privacy mechanism in a mechanism similar to BERG coins and tokens.

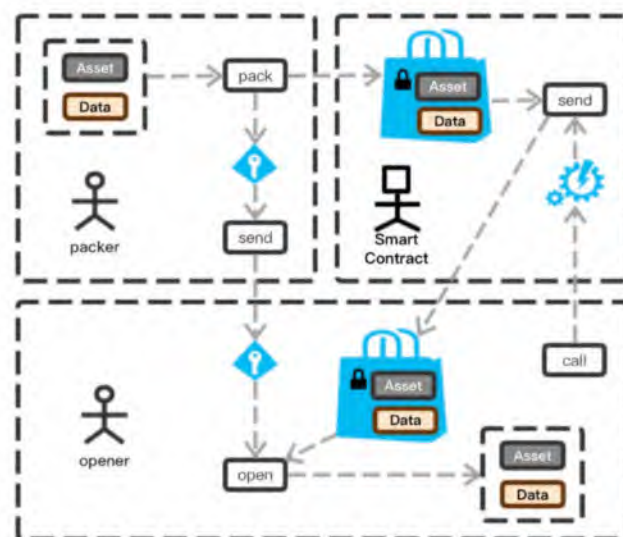


Ticket is a 256-bit digit that can point to complex data structures stored in smart contracts, so it is suitable for the construction of complex business scenarios. A concise example of the data structure pointed to ticket is shown in the figure above. *Asset* is the data structure of an asset. The data structure



of this asset can be BERG coin, token or ticket. Because Ticket is directional to data, this structure can support unlimited nesting and form a complex assets structure that meets business requirements. Data can store data encoded by other data structures other than assets, and the data structure represented by this data can also be nested. Package holds the assets or data encoded by the user with a key. With the feature of ticket, developers can completely transform complex game applications such as Crypto Kitties into anonymous versions.

#### H. Package: The Mechanism Supported by BERG System for Encrypted Assets or Packets



In data-centric applications, these data include timestamps, high-value data storage, various types of proofs data or password strings, asset data, etc. These applications, if implemented in blockchains and input and output by smart contracts, they will encounter exposure problems with plaintext data. BERG also provides the corresponding data protection technology, so that the exposure of these plaintext data is controlled by the user himself. This technology is called Package. The current version of the package uses the symmetrical encryption technology ChaCha20, which is fast and secure. The user can package the Package on the client and get the corresponding key. After packaging, the Package can be transferred anonymously or entered into the smart contract to perform the corresponding logical operation. In the process, the contents of the encrypted package are uncrackable.

When decryption is required, the user can privately transfer the key to the associated party according to their own situation, so that the related parties can use the key to decrypt related assets and data.

The four types of assets, BERG coins, Tokens, Tickets, and Packages, together constitute the ecosystem of the BERG system's anonymous assets. These asset types are perfectly integrated in a single coding system, combined with the programming of these assets by Turing's complete smart



contract of BERG system, developers can pioneer the implementation of various privacy-related DApps that are not supported on the previous system in the blockchain, and are suitable for the implementation of business requirements in various privacy protection fields.

### **3.6 ORIENTED SCENARIOS**

Privacy protection is a strong demand for individuals and organizations in the real world. BERG supports Turing's complete smart contracts and various related privacy components, which can support the expansion of different economic ecology. The issuing of the anonymous asset is no longer exclusive to a few that have extensive knowledge of cryptography. General system developers, with the business needs to issue anonymous cryptocurrency, can issue their own anonymous tokens on the BERG chain and establish their own privacy ecology, therefore, greatly expanding the application of the blockchain systems. Here are some typical use scenarios below.

#### **A. Supply Chain System**

The blockchain can solve the problem of upstream and downstream transaction vouchers and traceability of the supply chain system, simplify the management of supply chain central enterprises and providing corresponding solutions for financing of upstream and downstream enterprises. However, sensitive data such as prices and quantity of goods are facing the problem of leaking trade secrets when they are on the blockchain. With BERG system, the problem of exposure to trade secrets can be completely solved, and at the same time, the participating parties can enjoy the benefits brought by the application of blockchain system.

#### **B. Medical Health**

Digital privacy is in all aspects of the medical and health industry. From personal medical records to medical treatment records, the multi-role privacy protection and the authorization mechanisms require highly flexible and secure privacy protection capabilities, including hospitals, patients, insurance companies, pharmaceutical companies, etc. The BERG system solves the privacy problems faced by patients and hospitals, and also opens the way for insurance companies and pharmaceutical companies to be safely compliant and use relevant data with the permission of patients.

#### **C. Online Auction**

Online auction businesses pursue fairness, the privacy of bids is an important aspect and often difficult to obtain because of conflicting interests. BERG can provide a completely safe, independent and fair bidding environment.

#### **D. Online Casino**

The development of online casino industry has always been restricted by the centralization mechanism. The applications of Online casino often require a high degree of privacy protection for competitors' strategy. In these huge cash flow applications, the decentralized smart contract systems that can provide multiple bids, payment and settlement is needed, and BERG system can fully support this kind of business.

#### E. Online Gaming

Large-scale online games often need token systems that are easy to circulate, trade and settle, and can be issued and circulate based on smart contracts, while providing privacy protection of transactions. BERG is the only technical solution that supports a multi-token system that issues and circulates homomorphic smart contracts, with added transaction privacy for the accounts.

#### F. Conclusions

There are more industries involved in asset digitization and privacy-sensitive digital assets, such as insurance industry, commodities trading, futures trading, credit industry, digital asset trading (such as credit investigation and intellectual property rights, etc.) and so on. In these areas, BERG system has broad application prospects.

### **3.7 FUTURE PLANS**

#### **Off-Chain Computing and Homomorphic Encryption Smart Contracts**

The homomorphic encryption of smart contracts has already entered development stage and is planned to be released on BERG platform of version 2.0 in March 2019. The team discovered a method to balance data security (a mechanism that completely isolate sensitive data for the computations) and performance through both on-chain and off-chain computing. The plan aims to finish the work within 6 months.

#### **Wallets and Other Ecological Applications**

BERG's decentralized wallet application is currently under development and is scheduled to be released in March 2019. BERG supports developers to issue tokens themselves, the wallet will support BERG's own tokens and the management of cryptocurrency assets corresponding to all developers-based tokens issued by BERG.

#### **Latest Consensus Mechanism**

Within one year, the team will release a new consensus mechanism SE-Random in an updated version of BERG. The design will combine the latest PBFT theory and VRF algorithm in the consensus mechanism capable of balance fairness and efficiency.

#### **Privacy Three Swordsmen**

BERG has two related protocols, the Alien Protocol and the Castro Protocol. The former provides a distributed DNS system to obtain the stable operation of the network and information transmission by means of automatic addressing. The latter implements encrypted privacy protection for the IP address of the node, forming a complete decentralized application privacy protection scheme in the 3 in 1 suite.

### **Secure Multiparty Computing**

In many cases, data certification must combine with existing centralized data sources and can also become offline data sources. The current strategy to solve the above problems is to assume a trusted service provider or a trusted third-party exists. The risk is high in the changeable and malicious environment. A universal secure multi-party computation solution can resolve the problem.

BERG plans to introduce Secure Multi-party Computing (SMC) in the future, in order to provide extensive support of off-chain data under the premise of privacy protection.

### **Multi-chain system**

The multi-chain system is the BERG's scalability solution. BERG will use a mechanism similar to the Ethereum's Plasma for performance expansion based on multi-chain system, BERG's status updates per second can reach extremely high levels (possibly billions). This solution allows BERG to have the capability to replace today's centralized clusters with better performance, giving BERG the prospect of handling privacy-related decentralized applications around the world.



**BERG**

## **CHAPTER IV**

# **CHAIN FRAMEWORK**

In addition to the privacy protection mechanism, the chain infrastructure also needs to have sufficient scalability, an important aspect for building a practical application platform. BERG introduces the following technologies to enhance the chain's underlying architecture:

- Consensus Optimization - BERG uses a brand-new consensus mechanism SE - Random, that combines the latest PBFT theory and VRF algorithm to form a consensus mechanism balancing fairness and efficiency.
- Plasma - a method to obtain the expanded computing of blockchain. In Plasma, many blockchains are combined into a tree structure to participate in the computing, thus obtaining the horizontal expansion of the blockchain.
- Powerful Virtual Machines – the virtual machines not only meet EVM compatibility, but are also sufficiently scalable, and have the underlying instructions to meet performance requirements.
- The following focuses on some of the specific implementation of technologies.

### **4.1 CONSENSUS MECHANISM**

BERG proposes proprietary developed main chain consensus engine SE-Random based on the study of various consensus. The design of SE-Random Consensus Engine is influenced by the latest consensus research of Algorand and Ourboros. The verifications at the nodes have little computational overhead, and the probability of forks of the whole blockchain network is extremely limited. The engine can achieve almost unlimited scalability.

Detailed description of SE-Random consensus:

SE-Random uses Byzantine Agreement (BA\*) to reach the consensus in a set of transactions. For scalability, SE-Random uses a random algorithm to select a group of users, and allows users to check

privately whether they are selected or not to participate to form the consensus in the BA\* protocol. With the algorithm, as the number of users increases, the BA\* consensus system will not slow down.

### **The Use of VRF Algorithm**

SE-Random consensus engine is based on the Verifiable Random Function (VRF) algorithm for random verification node selection. VRF is a randomly generated function. The function is verifiable: the same private key signs the same information and only one legal signature can be verified. The function is different from the common asymmetric encryption algorithm.

The specific operation flow of VRF:

1. The prover generates a pair of keys, *PUB\_KEY* and *PRI\_KEY*. *PRI\_KEY* is the private key and *PUB\_KEY* is the paired public key.
2. The prover outputs random result:  $result = VRF\_Hash(PRI\_KEY, info)$
3. The prover outputs random proof:  $proof = VRF\_Proof(PRI\_KEY, info)$
4. The prover submits the random result and the random proof to the verifier. The verifier verifies whether the result and proof match, if matches, then proceeds to the next step
5. The prover submits *PUB\_KEY* and *info* to the verifier, and the verifier calculates whether the  $VRF\_Verify(PUB\_KEY, info, proof)$  result is *TRUE* or not.  $\forall TRUE$ , the verifier will pass the verification.
6. After verification, it can conclude whether *info* and *result* match; the data given by the verifier is correct. The verifier does not have the prover's private key *PRI\_KEY* during the process.

### **Random Seed Generation**

Seed will be used in random algorithm of some areas in SE-Random: the need of randomly selected and publicly disclosed seed in the encryption lottery of SE-Random. The seed must be known to the participating nodes but cannot be controlled by opponents. The seed generated by SE-Random in round *r* is determined by the seed of the previous round *r-1* by VRF. The seed and the corresponding VRF certificate are included in each proposed block. Once SE-Random agrees on the block of *r-1*, everyone will know the pseudo-random seed of the current round after the *r* round starts. The initial seed *O* value is calculated together by the initial participants using multiple nodes, resulting in an absolute unpredictable random seed. The resulting seed cannot be predicted or manipulated by attackers.

### **Method for Selecting Verifier by Encryption of Lottery through VRF Algorithm**

Se-Random uses an encrypted lottery method to select a random subset of users based on the weight of each user. The system sets a fixed number of BERG tokens as a screened candidate unit *S*, and specifies each node has a limited number of BERG coins *W* as the screening computation. The

total weight of all candidate units is  $W = \sum_i w_i$ . If node  $i$  has a BERG token of  $j$  screening units, node

$i$  can participate in the lottery screening as  $j$  different child nodes. The randomness of the cryptographic sortition algorithm comes from the random seed discussed in the previous section. SE-Random builds a VRF based on the current seed in each cycle of BA\*. The private key of VRF is only known by the node itself; each node uses its own private key to run a random algorithm published by the system. The system selects the verification node based on thBERGde holds no more than the defined proportion threshold of BERG coins.

SE-Random specifies a threshold to select the expected number of verification nodes. The expected number satisfies the probability  $p = w/W$ . The probability of a child verification node selection in  $W$  (total node weight) satisfies the binomial distribution:

$$B(A;IT,p) = \binom{W}{k} p^k (1-p)^{W-k} = 1$$

To determine the number of current verification nodes (including child verification nodes) is also determined by the cryptographic sortition algorithm. The cryptographic sortition algorithm divides the interval  $[0,1]$  into continuous intervals:

$$f = \left[ \sum_{k=0}^j (k-p)^{\wedge(k;w,p)} \right] \text{ for } j \in \{0,1, \dots, W\}$$

If the bit length of the hash is  $hashlen$  and if  $hash1 \cdot 2^{\wedge hashlen}$  is between  $j$ , then the node has  $j$  selected verification child nodes. The number of verification nodes selected can use  $u$  for VRF public authentication.

The characteristics of the cryptographic sortition:

1. The verification nodes randomly select  $N$  verification sub-nodes according to the weight of BERG tokens they hold.
2. Attacker that does not know the private key of node  $i$  cannot know whether  $i$  is selected, and how many sub-verification nodes are selected.

#### **BA\* consensus calculation performs in randomly selected verifier nodes**

Verification nodes (including sub-verification nodes) know that they are selected in secret and can only prove their verifier qualifications by publishing credentials. For each selected node, use signed seed with its own private key and enter a hash function to obtain its own credentials. The property of the hash function dictates the certificate is a 256 length of random string, the certificates of different nodes are not the same, and the distribution of the certificate strings is uniform. Using the same method, additional candidate leader nodes are selected. The candidate leader nodes are arranged

according to the certificate's lexicographic order. The smallest index order candidate is selected as the leader node, that is, the leader node is generated through a random public selection of the candidate leader node set.

The verification node and the leader node participate in the operation of BA\* protocol. In each stage and step of BA\*, the node independently determines whether it is selected in the current step through private and non-interactive means. BA\* has a two-stage voting mechanism. In the first stage, the verification node performs reduction consensus on the selected candidate blocks and selects the candidate block with the most consensus verification. In the second stage, binary Byzantine agreement is carried out on the candidate blocks screened in the first stage. BA\* consensus needs to ensure that more than two-thirds of the honest nodes participate in consensus. If the condition is not met, multiple random selections are required, when more than two-thirds of impartial nodes participate in consensus at once, consensus is reached. The verification nodes of each step of BA\* consensus is specified by cryptographic sortition in parallel to speed up the accelerate the consensus confirmation.

### BA\* Consensus Protocol

Every step of BA\*, the temporary key of the current step is destroyed:

#### 1. Generate Block (Step1)

- 1) The node checks whether it is a leader node  $B'$ .
- 2) Generates the message of the first step  $m^i = (b', \text{esg}, (h(b')), \langle \cdot \rangle)$
- 3) Broadcasts  $B'$  and  $m^i$

$m^i$  is the message that node  $i$  broadcasting at the  $(r, 5)$  step;  $B'$  is the block generated by the node  $i$  in  $r$  round;  $\text{esg}$ , refers to signing information with current temporary key at  $r$ ;  $H$  is a hash calculation function;  $\langle \cdot \rangle$  refer to the signature  $\text{sig}_i(r, s, Q^{-1})$  of  $I$ , which is used to prove the existence of  $I$  in the verification node set.

#### 2. Reduction Consensus

This protocol transforms the problem of agreement on any block into an agreement on two values. The two values are the basis for the finalizing the hash of a specific block or the hash of an empty block divided into three steps. The three-step process will be described in detail in the technical yellow paper. If more than two-thirds of the messages  $(\text{esg}^k, \wedge)$  match, the specific block can be broadcasted, and if not, empty block is broadcasted. The message is then used for subsequent binary Byzantine Agreement.

#### 3. Binary Byzantine Agreement

The verification node verifies the values issued by the reduction consensus protocol. The binary Byzantine Agreement is a three-step cycle; the verification node will continuously check the historical values received to see if there are two kinds ending conditions: whether the total number of votes



reaches  $2/3$ , the number is the number of votes that the block is legal or the block is illegal. If the block is illegal, the consensus system will determine and generate an empty block. To prevent the occurrence of infinite loops, a maximum total number of loops  $m$  is set. If an end condition is not met after reaching  $m$ , the consensus system will temporarily generate a tentative consensus and form a final consensus in the subsequent process (the next few rounds) and confirms the earlier transactions.

SE-Random consensus will adapt to consensus decision in the case of weak synchronization network. Strong network synchronization will not cause block fork. If weak network synchronization occurs, a tentative consensus will be specified temporarily and the final consensus will be reached after the strong network synchronization is restored. SE-Random can prevent Sybil attacks, selfish mining attacks, Nothing-at-Stake attacks, long-range attacks, and other attacks. If the users of BERG chain spread to more than 100 million nodes, the SE-Random consensus can quickly reach a network-wide Byzantine Agreement consensus through VRF mechanism.

## **4.2 PHASED EVOLUTION OF CONSENSUS**

Although the mechanism of SE-Random has great advantages in TPS, if implemented directly, it will bring BERG coins centralization problems, early users or get too many incentives, so that the verification ability of large nodes will soon exceed that of ordinary PoS miners, which is not conducive to proliferation. Therefore, BERG adopts a step-by-step consensus evolution process.

### **Pure PoW Consensus Phase**

In the early stage of Beta version of BERG chain, PoW consensus was adopted to facilitate the initial diffusion of BERG consensus among miners. The core design idea of the PoW consensus is to propose an operation procedure for calculating the value of complexity. Users can calculate a hash function by workload proof and submit it to service subset for quick verification to ensure the fairness and security of data transaction. The Beta chain of BERG is mined by Ethash consensus. The PoW of Ethash is memory-hard. Its PoW calculation needs to select a fixed subset of resources depending on nonce value and block head. The data structure of the subset is directed acyclic graph DAG. Every 30,000 blocks in Ethash generates a DAG, and the client needs to pre-generate the DAG in order to start the mining work.

### **PoW+PoS Consensus Phase**

The first MainNet version of the BERG chain uses the PoW + PoS hybrid consensus. With the development of BERG, the power of PoW miners is at risk of centralization. After the license mechanism is removed from the MainNet version, there is a theoretical possibility of 51% attack. The attacker obtains more than 51% of the power of the whole network by investing funds, carries out double spend attack or denial of service attacks by refusing to include transaction records in the block. Therefore, it is necessary to introduce PoW and PoS in the Mainland version, so that users can vote with their currency rights to prevent large-scale power attacks, so that large-scale power can effectively serve the main chain of book security, rather than malicious control of the system solely for the sake of economic interests. If an attacker wants to carry out a malicious carries out double spend attack, he needs to control the computing power of PoW while grasping a large amount of currency



rights, which greatly increases the cost of attack and reduces the possibility of double spend attack. In addition, the rights holder can decide which transactions are included in the block, so as to resist the inclusion of transaction records in the block. This can resist the blackmail of malicious PoW miners who refuse to include transaction records in the block to carry out denial-of-service attacks.

Pow+Pos hybrid consensus mechanism, by stimulating coin holders as PoS nodes, will get more online nodes and more effective stability of location block topology network than POW nodes with hardware input. Under the PoW + PoS mechanism, PoW is responsible for producing blocks, while PoS decides the legitimacy of blocks through the voting mechanism of rights and interests. In this way, the formation of each block is jointly completed by miners and money holders, and the two balance each other, thus avoiding the monopoly of either side. In addition, the voting mechanism of PoS can effectively suppress the hard bifurcation of block networks.

#### **SE-Random Consensus Phase**

The SE-Random Consensus is described in Chapter 4.1 above.

### **4.3 POW CONSENSUS PART OF MAINNET 1.X**

The PoW algorithm on Mainnet adopts ProgPoW. The purpose of using ProgPoW is to narrow the efficiency gap between commercial GPU and ASIC miners, to resist the computer monopoly of ASIC miners and to achieve fairer mining. ProgPoW algorithm is improved by Ethash algorithm, so the POW consensus algorithm of BERG BetaNet can switch to ProgPoW smoothly. Compared with Ethash, its resistance to ASIC is mainly manifested in the following characteristics:

1. Change keccak\_f!600 (64 words) to keccak\_f800 (32 words) to reduce its impact on the total calculation power.
2. Increase the mixing state
3. Adding Random Mathematical Sequences to Main Cycle
4. Add low-latency, small-scale cache reads that support random addresses
5. Increase DRAM reading from 128 bytes to 256 bytes

In addition, ProgPoW has six definable variables that can be bifurcated to further address threats from ASIC or FPGA. These six variables are as follows:

PROGPOW\_LANES: Number of parallel rows coordinated to compute a hash instance; default value is 32

PROGPOW\_REGS: Register file usage size; default value is 16

PROGPOW\_CACHE\_BYTES: Buffer capacity; default value is 16\*1024

PROGPOW\_CNT\_MEM: The number of frame buffer accesses defined as the external loop of the algorithm; the default value is 64

PROGPOW\_CNT\_CACHE: Number of cache accesses per cycle; default value is 8

PROGPOW\_CNT\_MATH: Number of mathematical operations per cycle; default value is 8

## **4.4 POS CONSENSUS PART OF MAINNET 1.X**

### **Overview of BERG PoS**

The main responsibility of BERG's POS consensus is to verify the validity of PoW's blocks. Users need to lock up a part of their BERG coins to bid for tickets. When a PoW miner generates a block, he needs to select three ballots randomly from the tickets pool to conduct a validated vote on the block. After the vote is completed, the BERG coins used by the user to bid for the ballot will be returned, and the corresponding block PoS reward will be obtained. In order to keep the number of PoS transactions within a reasonable range and to ensure that the validity of the validation is confirmed after the network blocks are further solidified, the PoS purchased ballots and block awards of BERG coins, which will be automatically settled and returned to the account approximately periodically after the users participate in the PoS.

When votes are voted on blocks, the holder of the ballots needs to keep the whole node online and complete the valid voting automatically by the account. If online voting is not maintained at that time, it will be treated as abstention and block award will not be obtained.

### **BERG PoS StakingNode**

When voting is needed when the votes are selected, users need to vote online through the nodes. If the votes cannot be voted, users can not be rewarded. Considering the cost of the nodes online and the network and hardware conditions needed to participate in voting validation in time, users can choose to vote by special nodes, which we call Staking Node.

StakingNode of BERG can be selected by any node actively through registration on the network. When registering StakingNode, a certain amount of BERG needs to be mortgaged to the network.

After successful registration as a StakingNode, the node can provide its own node address to ordinary PoS participants. When ordinary PoS participants participate in PoS, they can enter a valid StakingNode address to vote for them when submitting and purchasing votes, so that they can participate in PoS through StakingNode in order to get a bigger voting success rate.

All the StakingNodes in the whole network will share the block rewards. For those nodes like StakingNode, because they provide a stable voting channel and fairly distribute the service value of the PoS rewards, they will get a certain proportion of the block rewards.

### **BERG PoW+PoS Process**

1. PoS miners can buy voting tickets in the pool.

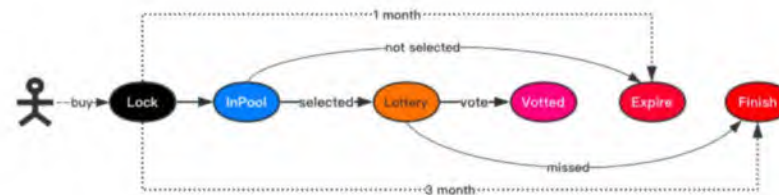
Shares have tickets pools, no upper limit, and the price is adjusted dynamically according to the size of the tickets pool. The change of share price is relatively balanced within the size of tickets pool 326592.

Share price rises sharply after the size of tickets pool exceeds 326592.

2. After calculating the block HASH, the PoW miners will randomly select three tickets in the tickets pool and broadcast the ballot information.

3. When the PoS miners receive their own ballot information, they immediately broadcast their own signed ballot information.
4. After receiving the voting information of PoS miners, the PoW miners put their signatures in the block head information.
5. If the correct votes exceed 2, the PoW miners will broadcast the block and complete the process.
6. PoS miners can vote SOLO or Staking Node instead.
7. Each block award is divided into PoW award and PoS award. The size of PoW and POS awards varies independently with their difficulty.
8. PoS award is divided into Solo voting and Staking Node voting. The ratio of Solo award and voting award is 3:4.
9. After the voting process is completed, the reward and principal involved will be returned in the latest payment cycle.

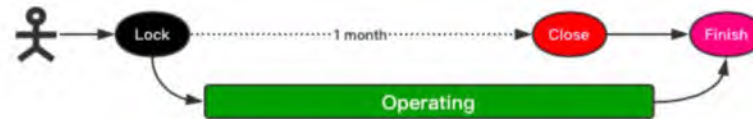
### Share's Life Cycle



1. After the user chooses a certain price to purchase the tickets, the principal will enter the Lock and the tickets will enter the InPool.
2. PoW miners randomly select three tickets when they exit the block, and the selected tickets are put into the state of waiting for voting (Lottery). At least two votes can be obtained to complete the block process.
3. When PoS miners receive the message waiting for voting, they successfully broadcast the voting information to PoW miners to complete the block. The tickets in the voted state will return the principal and reward within one payment cycle (one week) and release the tickets at the same time.
4. When there are only two successful votes in the block, the other share goes into Missed state. As a penalty, the tickets in the state of mistake will release the principal after three consecutive months of lock-in.

5. Shares that are not successfully selected in the pool will expire one month later and release the principal when they expire.

### StakingNode's Life Cycle



1. StakingNode operators acquire the right to operate StakingNode by locking a certain amount of BERG coins.
2. StakingNode can be closed manually only after one month of operation and enter Close state. New applications for tickets will not be allowed after closure, and operators are obliged to complete existing tickets in Staking Node.
3. StakingNode will return the principal and enter Finish state when both conditions are met at the same time. A. All the principal of the acting tickets are released, B. Enter Close state.

## 4.5 EXPANSION MECHANISM

Plasma is a framework for motivating and enforcing smart contracts execution. The framework can scale to a large number of state updates per second and supports a large number of decentralized financial applications worldwide on the blockchain. The smart contracts use network transaction fees to stimulate continuous automation and ultimately rely on the underlying blockchain to force the finalizing of the transaction status.

Plasma consists of two core parts: reorganizing all blockchains into a set of MapReduce functions, and an optional method to implement a POS token deposit mechanism based on Nakamoto consensus principle that does not encourage block withholding on existing blockchains.

The mechanism enforces state locking on the main chain by writing smart contracts on the main chain and using fraud proofs. Plasma groups the blockchains into a hierarchical tree structure, treats each blockchain as an independent branch and force the history of the entire blockchain and MapReduce calculations to be submitted to Merkle Proof. Through the main chain to force encapsulation of the account book information of a chain into a sub-blockchain. The chain will expand from the lowest trust to achieve an incredible expansion capacity.

Withholding attack on blockchain is a complicated problem surrounding the global availability of data that enforces non-global data. Plasma alleviates the problem by withdrawing the problematic chain and also creates a mechanism for encouraging and enforcing the accuracy of the execution data continuously.

By broadcasting the normal status of Merkle's proofs to the main chain (e.g. Ethereum), this will achieve large scalability, reduce transaction costs and computational effort. Plasma supports the continuous operation of large-scale decentralized applications. The important scalability is realized by reducing the single expense financial expression to one bit in a bitmap. A transaction and a signature represent a transaction aggregation with multiple parties. Plasma combines this feature with a MapReduce framework and uses smart contracts with deposits to build a scalable computing compulsions.

The structure of the mechanism allows external participants to hold funds and calculate contracts according to their own action, similar to a miner, but Plasma runs on an existing blockchain so that users do not need to create corresponding transactions on the main chain every time the state is updated (including the addition of a new user's account ledger), only a small amount of information is written to the main chain such as the merged state change.

BERG will use a mechanism like Plasma for horizontal performance scalability based on multi-chain systems. The multi-chain parallel computation system allows BERG to have the ability to update its state at a very high level per second. As a result, BERG will greatly improve its performance to replace the carrying capacity of the current centralized cluster.

## **4.6 VIRTUAL MACHINES**

At present, Ethereum has a large number of developers, and Solidity language has become the most widely used language for smart contract development. Therefore, EVM compatibility is provided in BERG systems.

EVM virtual machine was developed on the basis of Ethereum. Ethereum has a standard blockchain structure with a simplex data structure. The virtual machine is designed to resemble the ACID (Atomicity, Consistency, Isolation, Durability) feature of the database at the transaction invocation level. In Ethereum's protocol, the invocation of one smart contract may affect the status changes of multiple accounts. The state changes are rigid transactions with real-time consistency, the state changes either occur simultaneously or do not occur at all. BERG considers scalability in the future and has the underlying instructional basis to meet performance requirements. The virtual machine of BERG-Chain satisfies the concept of BASE (Basically Available, Soft State, Eventual Consistency), and the virtual machine is referred to as MEVM virtual machine.

In the BASE concept, the Basic Availability means that the system is allowed to lose part of its availability in case of unexpected failure. Soft State means allowing the data in the system to have an intermediate state, but the existence of the intermediate state will not affect the overall availability of the system. Eventual Consistency refers to all copies of data will eventually be reaching consistency after synchronization. Compared to the strong consistency of ACID concept, BASE concept gains

usability by sacrificing the real-time consistency and reaches a consistent state eventually. The principle of the block structure and various consensus algorithms in the blockchain are essentially the BASE concept, but they do not conform with ACID. MEVM are designed to combine with BASE semantics. Compared to the original ACID design of EVM, the design will overcome the performance bottleneck.

In addition, the Solidity language has been criticized for lack of support of standard libraries, such as the basic functions of comparing two strings. There are no standard library functions in Solidity for developers to invoke. Projects like OpenZeppelin provide some standard libraries, but they are far from sufficient. BERG's blockchain applications require advanced mathematical and cryptographic algorithm libraries, such as Zero Knowledge Proof Protocol, RSA Public-key encryption algorithms, singular value decomposition, etc. MSolidity can refer to the implementations and add more libraries, which can be precompiled or implemented in native mode to reduce operational overhead.

In the future, BERG system may support virtual machines based on Web Assembly (WASM) to further improve performance and provide support for smart contracts written in languages other than Solidity (C, C++, Rust, or GO). As the IELE virtual machine designed by the Cardano project team matures, Matter system will consider providing support for this virtual machine. IELE is a variant of LLVM and is expected to become a unified and low-level platform for smart contract translation and execution in high-level languages. With IELE virtual machine, BERG system can support more variety of advanced languages.

## **4.7 POST QUANTUM CRYPTOGRAPHY**

Blockchain system commonly use asymmetric cryptographic signature algorithm, such as the RSA algorithm based on large integer factorization problem and the ECC algorithm based on the discrete logarithm computation problem on the elliptic curve, that can be easily cracked by quantum Shor algorithm which change NP problems to P problems. The BERG system will iteratively introduce encryption algorithms to safeguard against quantum computing, such as Lattice-based cryptography, code based crypto-systems, and multivariate cryptography, based on project progress and the development of quantum computer utility. The various crypto-systems such as encryption, signature, and key exchange can be designed based on the lattice password, which is an important direction of the post quantum cryptography algorithm. At the same time, the BERG team will continue to track the cutting-edge research directions of anti-quantum crypto-systems such as the Isogen problem, the conjugacy search problem, and the Braid Groups.



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## **CHAPTER V**

# **ECONOMIC MODEL**

The traditional point-to-point communication network focuses on information transmission, a bit like the application of Internet 1.0. Things are open and shared, unlike the disruptive blockchain technology. Because all human behaviors are driven by the economic logic, human behavior in the absence of effective economic norms can only be bound by social norms (i.e. work driven by spiritual incentives of public interest), which lacks the binding needed for individuals to complete the goals together.

Bitcoin network through the POW consensus mechanism and the contribution of using computing power to obtain the bookkeeping rights to obtain the bitcoin rewards to encourage the node to participate in the consensus is undoubtedly a remarkable design. The token economic model is the core of the value of a blockchain.

However, the question is whether one type of token can solve the incentive problem of every consensus cooperative behaviors? We think the answer is NO. There are various kinds of tokens circulating in the market, and the economic models behind them are varied, but there is the lack of a unified standard that link the cost of consensus with the consensus value generated. Therefore, the secondary market circulation rules of the cryptocurrency system appear quite fragile.

BERGd on the same underlying consensus mechanism, Ethereum allows smart contract developers to issue their own token and use ETH as a GAS fee to pay for the consensus cost, which unifies the unit of measurement of consensus costs, and allows users to obtain different value outputs according to the token's ecosystem. The users can at least calculate the best balance between investment and return. Many in the industry criticize the issuing of ERC20 tokens on Ethereum as too simplistic and that it can result in fraud, but few critics realize the importance of Ethereum's original design.

BERG team extended the features of Ethereum when designing the BERG-Chain. First, to reduce the GAS consumption in order to reduce the hard threshold of price-performance ratio for transactions on the chain. The team have designed a new consensus mechanism, which is described in another chapter.

Assuming that the consensus cost is negligible, the value of any token depends on other costs of transactions on the chain; which are affected by the centralization of digital assets and the relationship between market supply and demand, etc. The characterization is similar to real-world currencies.



Cryptocurrencies can also be used to measure the value of goods, services or rights, and interests of goods. Therefore, the developers should have their own economic model to issue tokens. The discussion of the economic model is from the aspect of BERG token.

From BERG's ecology, the value of all goods and services has a source. The blockchain platform is essentially a fair-valued circulation market circulation. The underlying cost of all economic activities is the transaction cost, and BERG token becomes the carrier of transaction cost. From this perspective, BERG token will be used for the following incentive purposes:

- Bookkeeping rewards;
- Computational contribution rewards ( more computing power consumption will be required for applications which use privacy mechanisms );
- Other roles including operational rewards for algorithm providers (by issuing smart contracts)
- In SE-Random consensus, possession of BERG's token could impact some specific scenarios (such as random selection of initial seed nodes);
- Developers of BERG ecology will get token rewards from BERG based on the actual value of the development and application. The rewards could be given in the forms of subsidize the consensus bookkeeping cost or computational power contribution.
- Users can use BERG token for various purposes in their DAPP or BERG related ecosystems.





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## **CHAPTER VI ROAD MAP**

### **6.1 DRAGONS OF AUTUMN TWILIGHT (VO.X)**

#### **2018.9 Release AlphaNet Network**

- \* Open source to GitHub
- \* Support anonymous encryption of common transaction information
- \* Support Turing's complete smart contracts
- \* Support issuing anonymous tokens using smart contracts

#### **2018.11 Release BetaNet-RC Network**

- \* Release foil-node PC user-interface Wallet
- \* Support issuing anonymous tickets using smart contracts
- \* Support decentralized mining license
- \* Support block browser

#### **2018.12 Release BetaNet-Release Network**

- \* Support issuing encrypted package
- \* Support sealed bid and private OTC transactions using smart contracts

- \* Support paying gas on behalf in smart contracts

## 6.2 DRAGONS OF WINTER NIGHT (V1.X)

### 2019.3 Global Node Deployment, Preparing for Main Network Environment

- \* PoW Mine Pool System
- \* Anti-ASIC and FPGA miner ProgPoW workload proof mechanism and mining client
- \* Centralized Mobile Wallet
- \* Mobile Wallet Integrating with BERG and ECO-tokens Exchange
- \* Perfect Large Stock Exchange Docking Middleware

### 2019.7 Release MainNet Network

- \* PC Decentralized Light Wallet
- \* PoS+PoW Hybrid Consensus
- \* Decentralized PoS StakingNode

### 2019.9 Light wallet enables interaction with smart contract

## 6.3 DRAGONS OF SPRING DAWNING (V2.X)

### Starting from the fourth quarter of 2019: Layer2 work is performed

- \* Inter-Blockchain Middleware: Inter-Blockchain with ETH and implement the Inter-Blockchain and anonymous processing of the tokens on ETH in the early stage
- \* Instant Trading Middleware: Node-based Layer 2 instant trading function under the chain to improve user's experience and achieve higher TPS
- \* Eco-middleware: Some middleware related to eco-applications

## 6.4 DRAGONS OF SUMMER FLAME

- Release ALIEN Protocol and CASTROL Protocol
- Complete version of SE-Random consensus
- Partitioning, Plasma and other related TPS enhancement schemes
- Anti Quantum Computer Function
- Privacy Protection Mechanisms Related to Downlink Data



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## **APPENDIX**

### **A LEGAL STATEMENT**

The sale ("Token Sale") of BERG Token is only used as an exchange medium for specific targeted crowds or participants. This is not any form of prospectus or offer document, nor is it BERG constitute any form of securities offer, unit in a commercial trust, unit in a collective investment plan or any other form of investment, or any form of investment offer in any jurisdiction. No regulatory organization has reviewed or approved any of the information listed in this white paper. This white paper has not been registered with any regulatory authority in any jurisdiction. By accessing and / or accepting any information in possession of this white paper or part thereof, as the case may be, by default you meet the following conditions:

( a ) You are not in the People's Republic of China, nor are you a citizen or resident ( tax or otherwise ) of the People's Republic of China, or reside in the People's Republic of China;

( b ) You are not in the United States of America, nor are you a citizen, resident ( tax or otherwise ) or green card holder of the United States of America, or reside in the United States;

( c ) According to the laws, regulations or rules of your region, you are not in a jurisdiction that prohibits, restricts or unauthorized sale of tokens in any form or manner, whether in whole or in part;

( d ) You agree to meet the conditions and constraints described above.

### **B RISK INDICATION**

This information does not represent an investment proposal, or any license for sale, or to guide and attract any purchase.